

# Automatically Proving Memory Safety and Termination of C-Programs

Alexander Weinert  
RWTH Aachen University

Research Area Computer Science 2  
Published: (Termination with Pointer Arithmetic, Ströder et al., 2014)

Feb 23, 2015

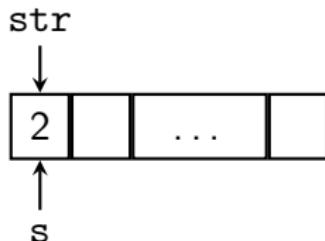
## An Example

```
int strlen(char* str) {  
    char* s = str;  
    while(*(++s));  
    return s-str;  
}
```



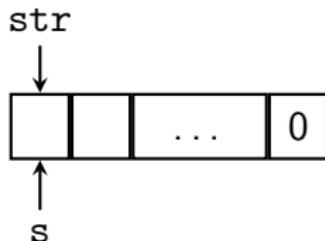
## An Example

```
int strlen(char* str) {  
    char* s = str;  
    while(*s++);  
    return s-str;  
}
```



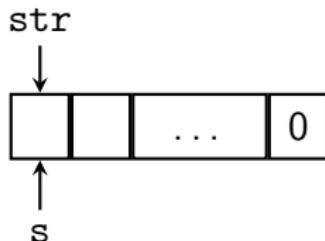
## An Example

```
int strlen(char* str) {  
    char* s = str;  
    while(*(s++));  
    return s-str;  
}
```



## An Example

```
int strlen(char* str) {  
    char* s = str;  
    while(*s) s++;  
    return s-str;  
}
```

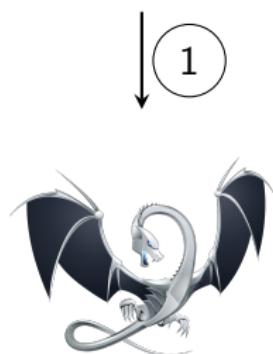


# Big Picture

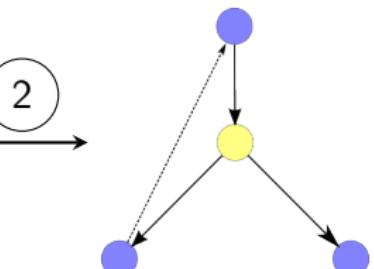
THE  
**C**  
PROGRAMMING  
LANGUAGE



YES/NO/  
MAYBE



LLVM



$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$

Integer Transition  
System

4



# C to LLVM

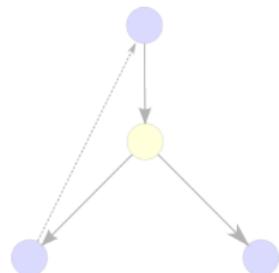
THE  
**C**  
PROGRAMMING  
LANGUAGE



LLVM



Symbolic Execution Graph



YES/NO/  
MAYBE

?

$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$

Integer Transition  
System

# LLVM Compiler Infrastructure

```
int strlen(char *str) {  
    char *s = str;  
    ...
```

C-program

```
define i32 strlen(i8* str) {  
    c0 = load i8* str  
    ...
```

LLVM Internal  
Representation

```
strlen_entry:  
    push edi  
    ...
```

Assembler

# Some LLVM Instructions

- ▶ Control Flow Instructions
  - ▶ ret, br, call, ...
- ▶ Arithmetic Instructions
  - ▶ add, icmp, ...
- ▶ Bitwise Instructions
  - ▶ shl, and, ...
- ▶ Memory Instructions
  - ▶ alloca, load, store, ...

Complete Reference: <http://llvm.org/docs/LangRef.html>

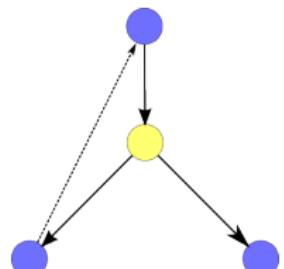
# LLVM to Symbolic Execution Graph



LLVM



Symbolic Execution Graph



$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$

Integer Transition  
System

YES/NO/  
MAYBE

?

## Problem Definition

Given: LLVM Program, Entry Point

Goal: Description of *at least all possible runs* from the entry point

Idea: Abstract interpretation of program states

# Abstract States

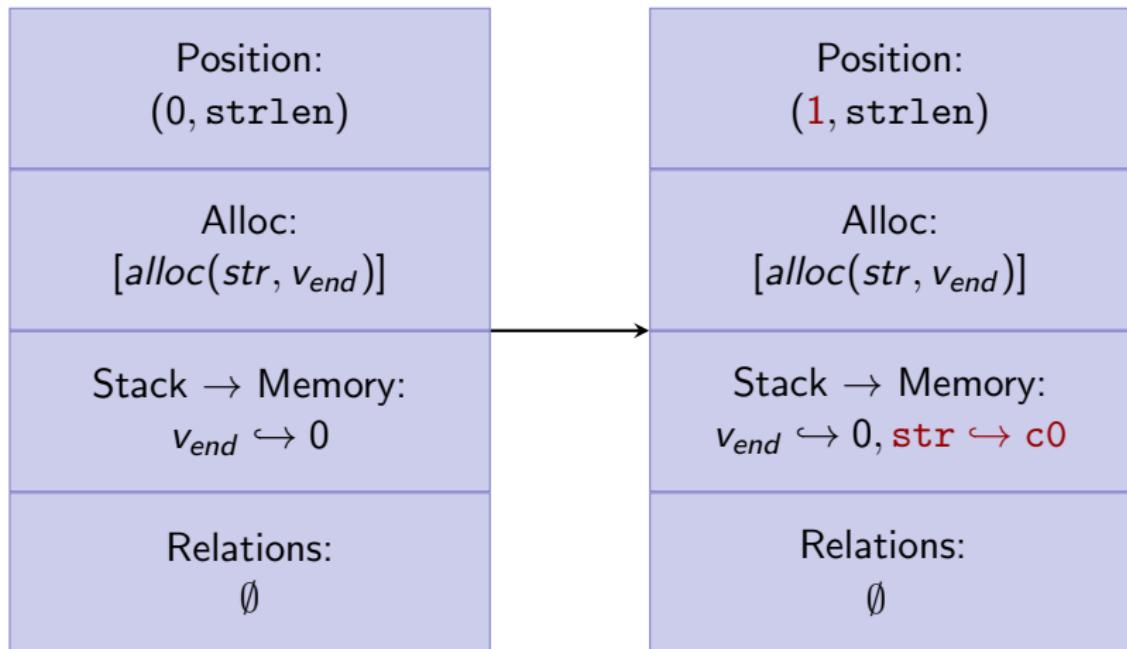
Position:	(0, $\text{strlen}$ )
Alloc:	$[\text{alloc}(\text{str}, v_{end})]$
Stack → Memory:	$v_{end} \hookrightarrow 0$
Relations:	$\emptyset$

Entry Point:

```
define i32 strlen(i8* str) {  
    c0 = load i8* str  
    ...  
}
```

# Evaluation of Abstract States

c0 = load i8\* str



# Refinement of Abstract States

$$c0zero = \underbrace{\text{icmp eq i8 c0, } 0}_{c0==0}$$

Position: (1, <code>strlen</code> )
Alloc: $[alloc(str, v_{end})]$
Stack → Memory: $v_{end} \hookrightarrow 0, str \hookrightarrow c0$
Relations: $\emptyset$

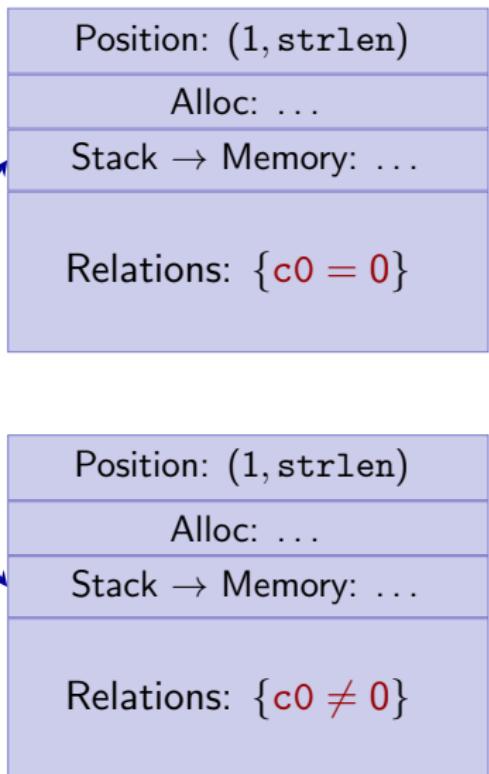
Position: (2, <code>strlen</code> )
Alloc: $[alloc(str, v_{end})]$
Stack → Memory: $v_{end} \hookrightarrow 0, str \hookrightarrow c0$
Relations: ?



# Refinement of Abstract States

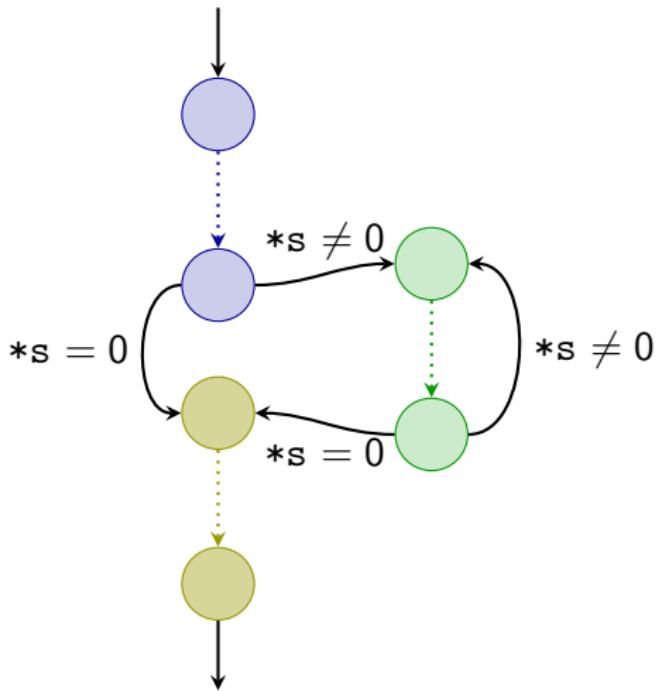
c0zero = icmp eq i8 c0, 0

Position: (1, strlen)
Alloc: [alloc(str, v <sub>end</sub> )]
Stack → Memory: $v_{end} \hookleftarrow 0, str \hookleftarrow c0$
Relations: $\emptyset$

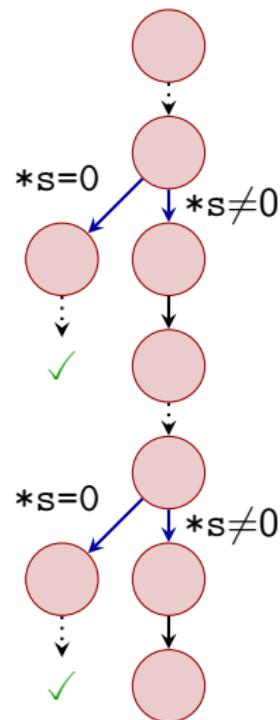
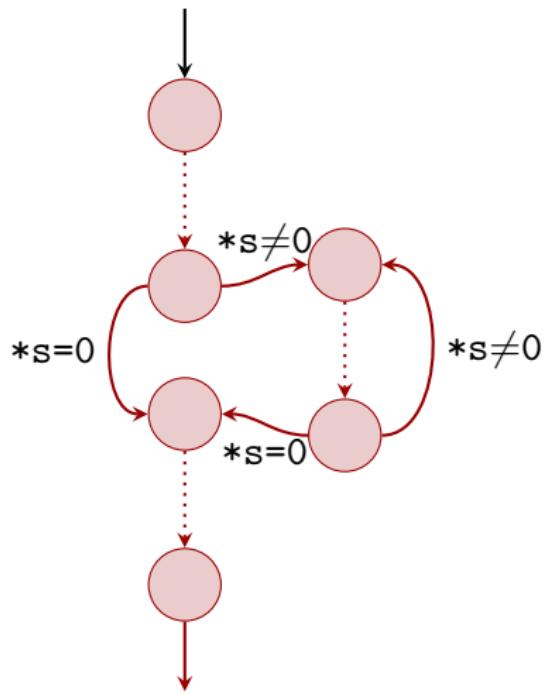


## Merging of Abstract States

```
int strlen(char *str) {  
    char *s = str;  
    while(*s) s++;  
    return s-str;  
}
```



## Merging of Abstract States



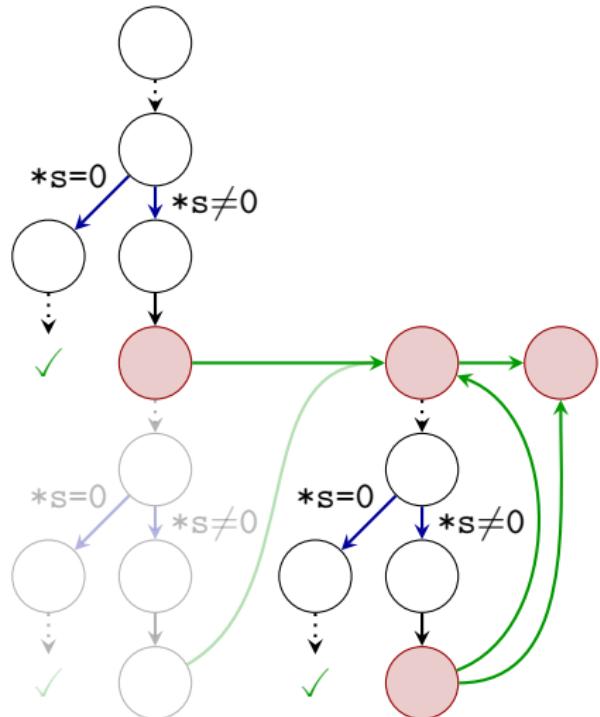
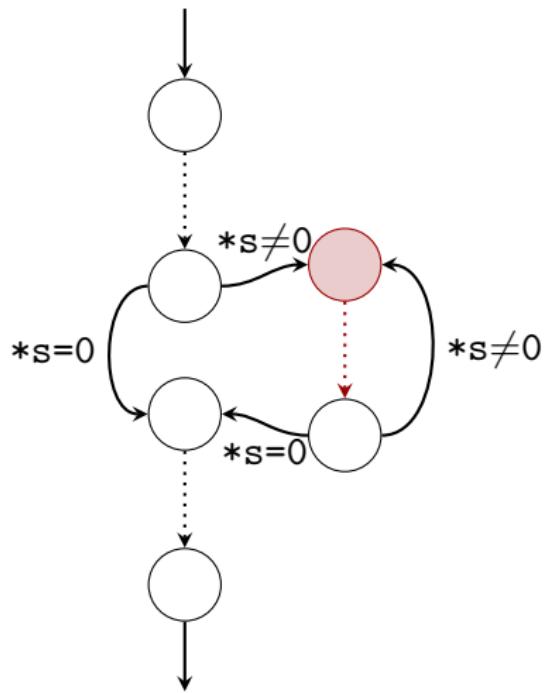
# Merging of Abstract States

Position: (3, <code>strlen</code> )
Alloc: [ $\text{alloc}(\text{str}, v_{end})$ ]
Stack → Memory: $s \hookrightarrow_{i8} c, \dots$
Relations: $\{s = str + 1, \dots\}$

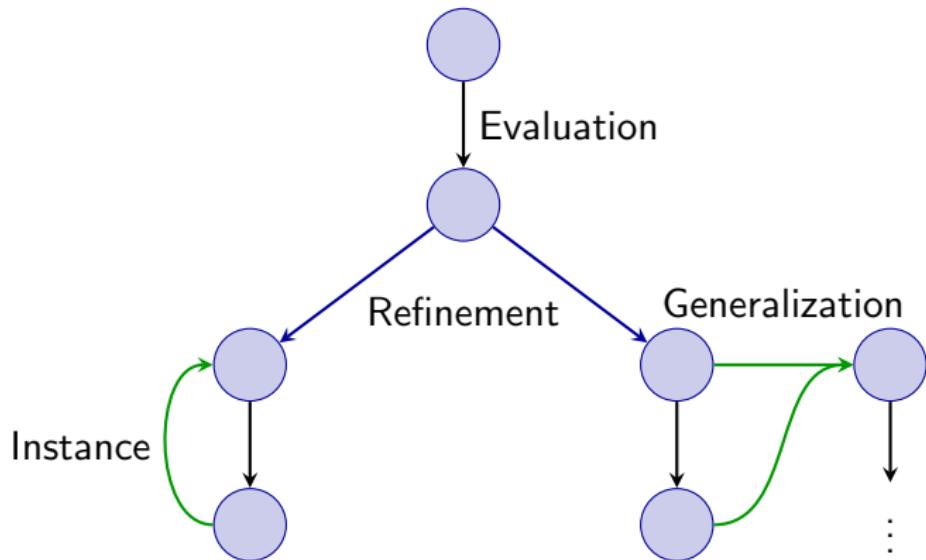
Position: (3, <code>strlen</code> )
Alloc: [ $\text{alloc}(\text{str}, v_{end})$ ]
Stack → Memory: $s \hookrightarrow_{i8} c, \dots$
Relations: $\{s = str + 2, \dots\}$

Position: (3, <code>strlen</code> )
Alloc: [ $\text{alloc}(\text{str}, v_{end})$ ]
Stack → Memory: $s \hookrightarrow_{i8} c, \dots$
Relations: $\{s > str, \dots\}$

## Merging of Abstract States



## Summary: Symbolic Execution Graphs



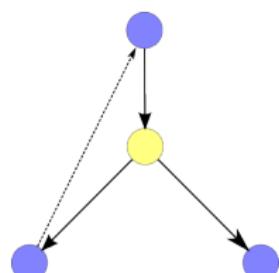
# Symbolic Execution Graph to Integer Transition System

THE  
**C**  
PROGRAMMING  
LANGUAGE



LLVM

Symbolic Execution Graph



YES/NO/  
MAYBE



$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$

Integer Transition  
System

# Integer Transition Systems

Term Rewriting System:

$$f(x, y) \rightarrow g(y, x)$$

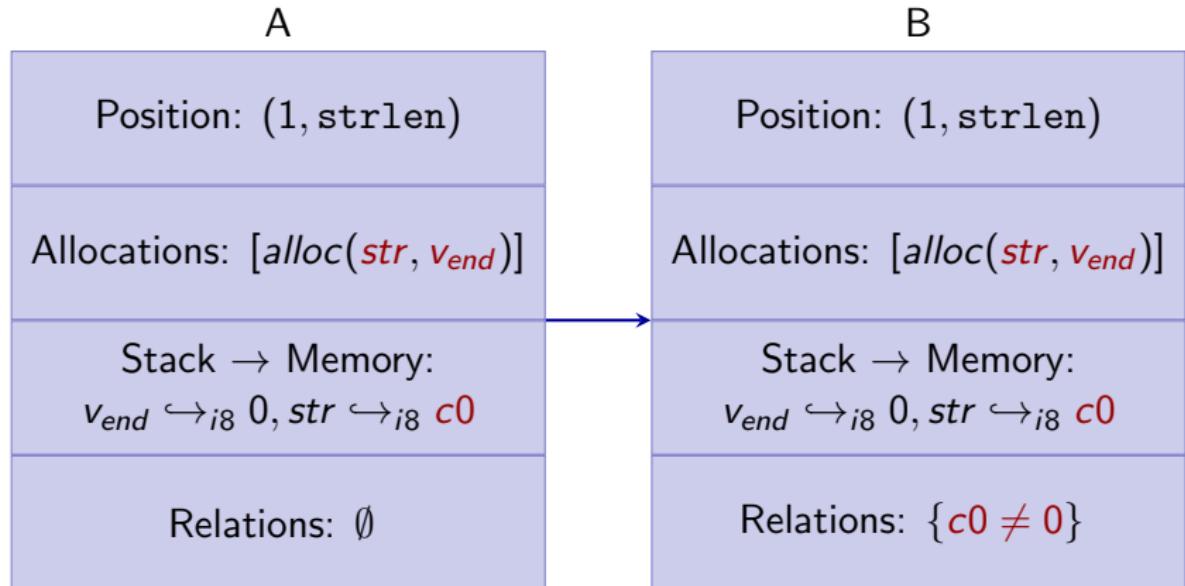
Integer Transition System:

$$f(x, y) \rightarrow g(y + 1, x - 2)$$

Integer Transition System:

$$f(x, y) \rightarrow g(y + 1, x - 2) \quad | \quad x > 0$$

# Symbolic Execution Graph to Integer Transition System



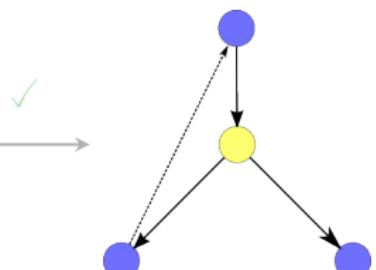
$$f_A(\textit{str}, \textit{v}_{\text{end}}, \textit{c}0) \rightarrow f_B(\textit{str}, \textit{v}_{\text{end}}, \textit{c}0) \mid \textit{c}0 \neq 0$$

# Termination of Integer Transition System

THE  
**C**  
PROGRAMMING  
LANGUAGE



LLVM



Symbolic Execution Graph

YES/NO/  
MAYBE

?

$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$

Integer Transition  
System

# Termination of Integer Transition System

Well-studied problem

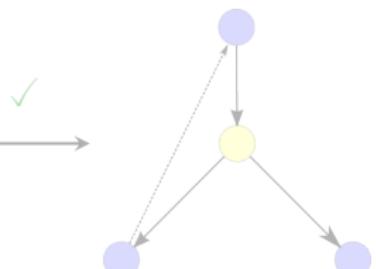
Use known techniques to show termination,  
e.g. (Termination of Integer Term Rewriting, Fuhs et al., 2009)

# Termination of Integer Transition System

THE  
**C**  
PROGRAMMING  
LANGUAGE



LLVM



Symbolic Execution Graph

YES/NO/  
MAYBE



$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$

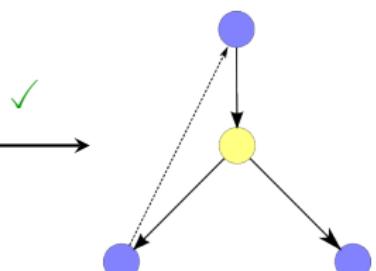
Integer Transition  
System

# Overview

THE  
**C**  
PROGRAMMING  
LANGUAGE



LLVM



Symbolic Execution Graph

YES/NO/  
MAYBE

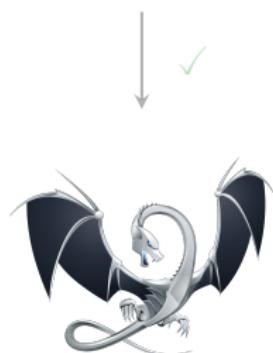


$$f(x) \rightarrow g(x + 1)$$
$$g(x) \rightarrow g(x - 1)$$

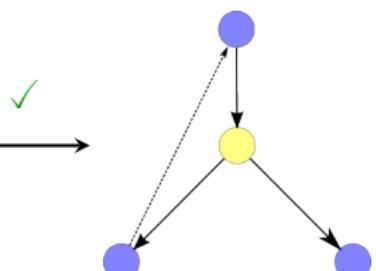
Integer Transition  
System

# My Contributions

THE  
**C**  
PROGRAMMING  
LANGUAGE



LLVM



Symbolic Execution Graph

YES/NO/  
MAYBE



$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$

Integer Transition  
System

## My Contributions

Given: Abstract State  $s$ , Integer Relation  $r$

Question: Does  $s \models r$ ?

## My Contributions

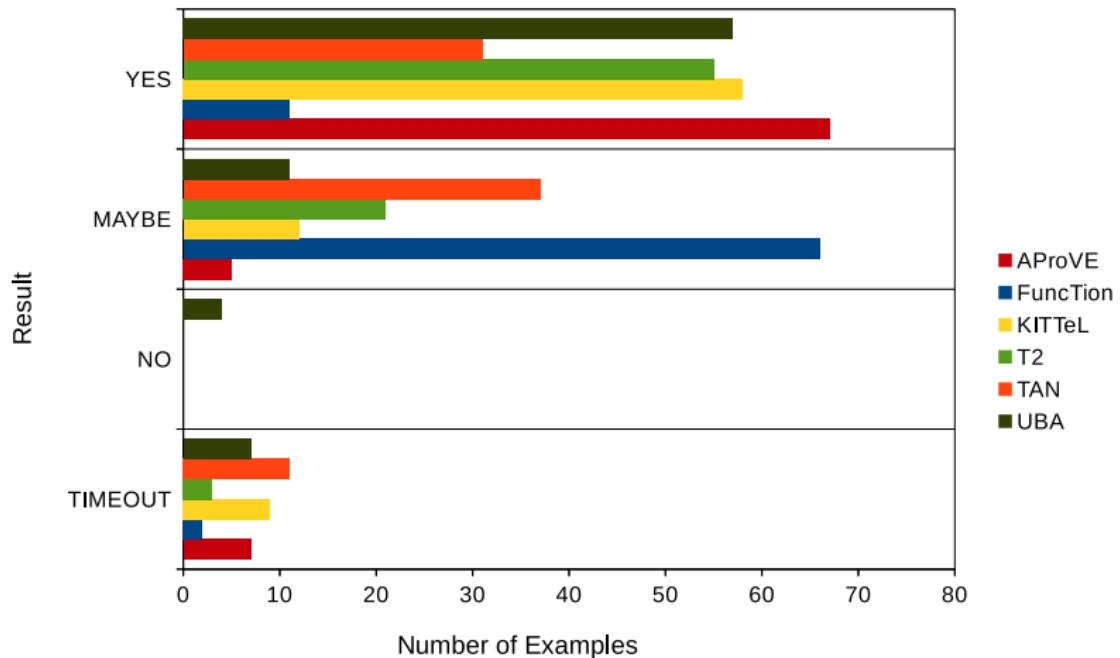
- ▶ Reduction from state to set of arithmetic relations
- ▶ Inference of knowledge from states
  - ▶ Formulation of inference in terms of integer relations
  - ▶ New framework for inference
  - ▶ Parameterization of framework with abstract arithmetic domain
  - ▶ Formulation of existing inference in framework

# My Contributions

- ▶ Use of Octagon Domain for inference of relations
  - ▶ (The Octagon Abstract Domain, Miné, 2006)
- ▶ Experimental comparison
  - ▶ “Traditional” inference
  - ▶ Inference in framework
  - ▶ New inference based on Octagons

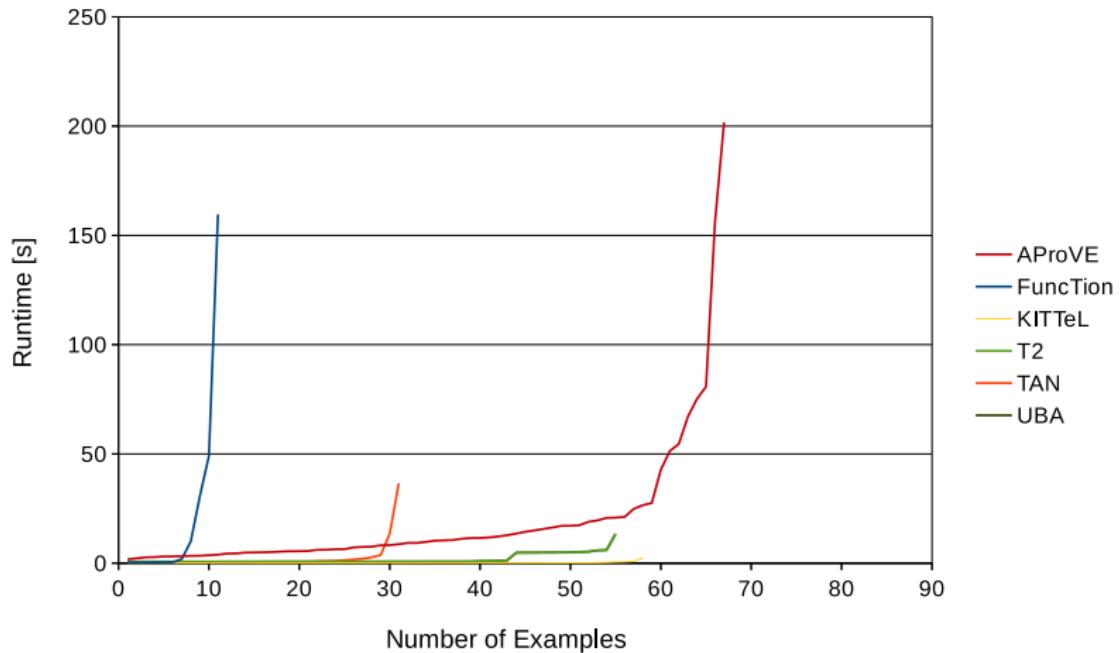
# Evaluation: Empirical Evaluation

79 Integer Programs, Timeout: 300s



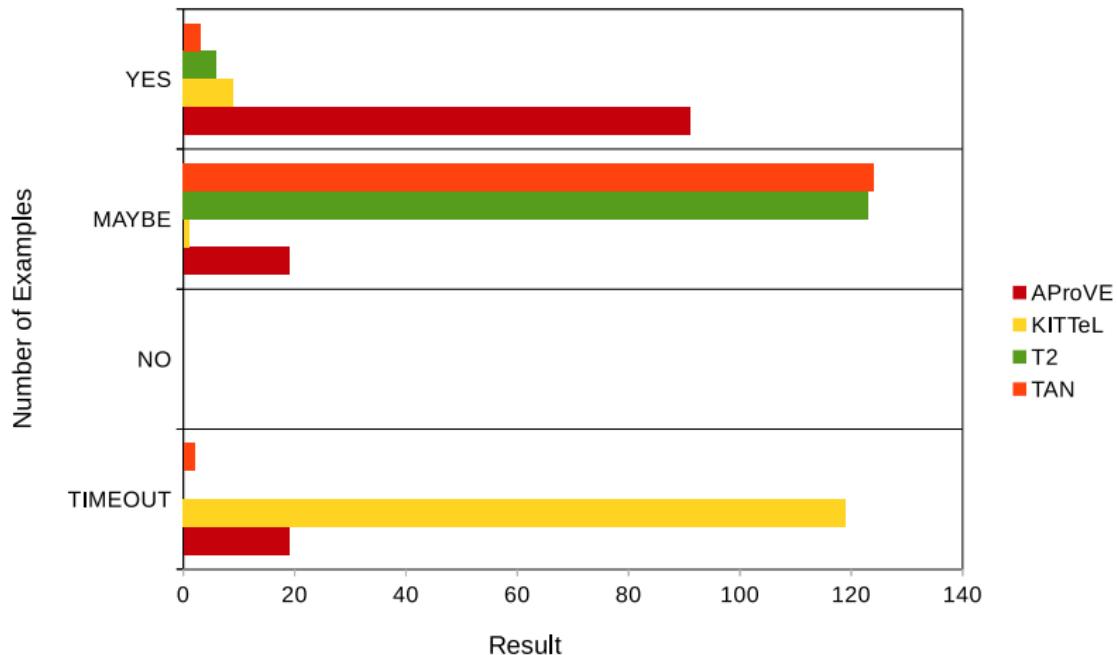
# Evaluation: Empirical Evaluation

79 Integer Programs, Timeout: 300s



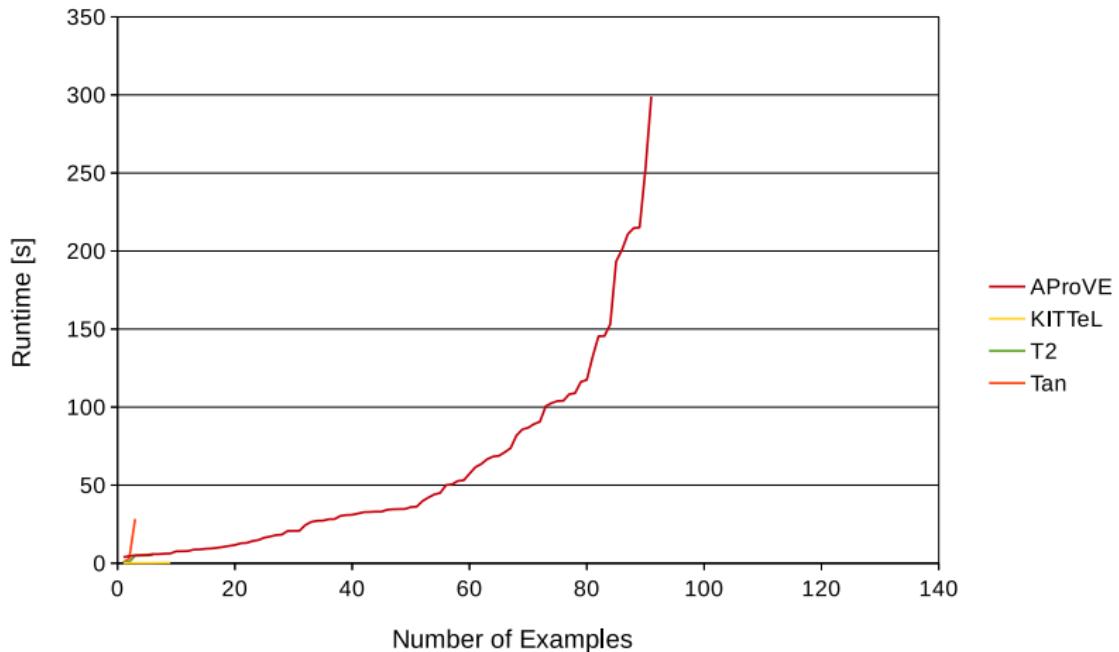
# Evaluation: Empirical Evaluation

129 Pointer Programs, Timeout: 300s



# Evaluation: Empirical Evaluation

129 Pointer Programs, Timeout: 300s

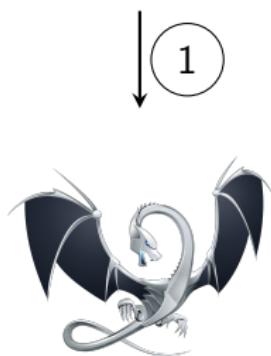


Complete evaluation at  
<http://aprove.informatik.rwth-aachen.de/eval/Pointer/>

# Thank you for your attention

[www.alexanderweinert.net](http://www.alexanderweinert.net)  
alexander.weinert@rwth-aachen.de

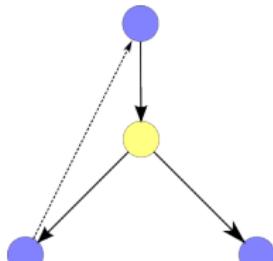
THE  
**C**  
PROGRAMMING  
LANGUAGE



LLVM



Symbolic Execution Graph



Integer Transition  
System

YES/NO/  
MAYBE



$f(x) \rightarrow g(x + 1)$   
 $g(x) \rightarrow g(x - 1)$