Analysis of Arithmetic Prolog Programs using Abstract Interpretation

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Master Thesis Presentation

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Roadmap



Question: Do all evaluations of the program terminate?

Prolog - Introduction

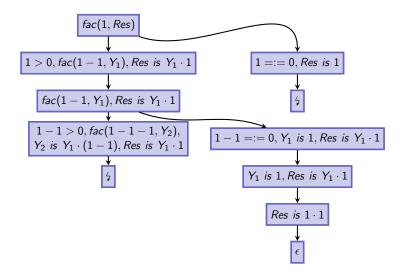
```
fac(X) =
    if X > 0 then Y1 = fac(X-1), return Y1 * X
    if X == 0 then return 1

fac(X, Y) :- X > 0, fac(X - 1, Y1), Y is Y1 * X.
```

fac(X, Y) :- X =:= 0, Y is 1.

PROLOG - Evaluation

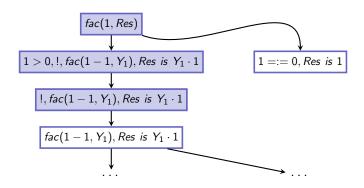
```
fac(X, Y) :- X > 0, fac(X - 1, Y1), Y is Y1 * X.
fac(X, Y) :- X =:= 0, Y is 1.
```



PROLOG - Cut

```
fac(X, Y) :- X > 0, !, fac(X - 1, Y1), Y is Y1 * X. fac(X, Y) :- X =:= 0, Y is 1.
```

PROLOG - Cut

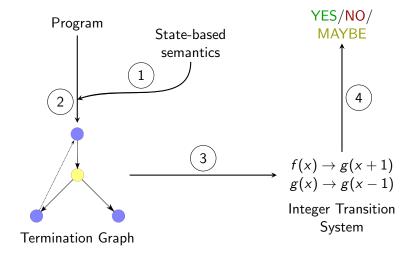


Termination

Given: $\ensuremath{\mathrm{PROLOG}}$ Program, some query template

Question: For all queries matching the template: Does the inference of the query on the program eventually terminate?

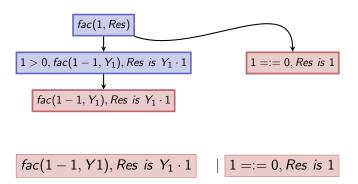
Roadmap



- ▶ PROLOG: Tree-based semantics
- Well-known techniques: State-based semantics

Solution: State-based semantics for PROLOG (Linear Operational Semantics, Ströder et al., 2012)

- ▶ Basic Idea: Leaves of tree describe state of inference
- ▶ Use front of tree as state



```
fac(X, Y) :- X > 0, !, fac(X - 1, Y1), Y is Y1 * X.
fac(X, Y) :- X = 0, Y is 1.
fac(1, Res)\downarrow1 > 0, !, fac(1 - 1, Y_1), Res is Y_1 \cdot 1 | 1 = 0, Res is 1
```

$$1>0 \ ,!, \textit{fac}(1-1, Y_1), \textit{Res is } Y_1 \cdot 1 \mid 1 =:= 0, \textit{Res is } 1$$

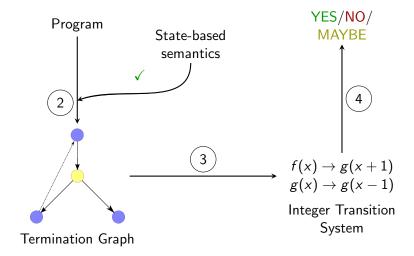
$$\downarrow$$

$$!, \textit{fac}(1-1, Y_1), \textit{Res is } Y_1 \cdot 1 \mid 1 =:= 0, \textit{Res is } 1$$

$$\downarrow$$

$$\textit{fac}(1-1, Y_1), \textit{Res is } Y_1 \cdot 1$$

Roadmap



From Programs to Graphs

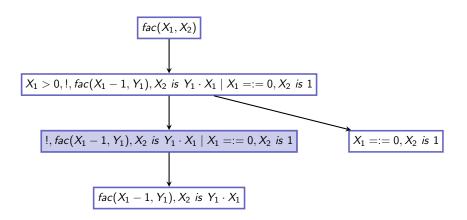
Given: Some Program, some query template

Goal: Finite representation of all possible inferences

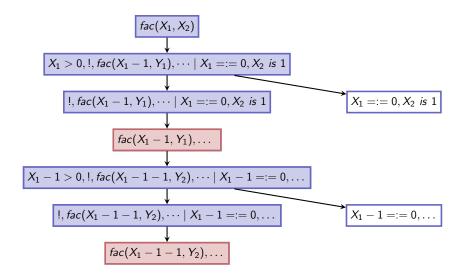
Idea: Represent set of runs as graph

From Programs To Graphs - Starting State

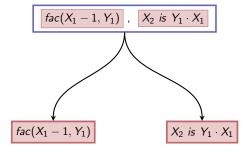
```
fac(X, Y) :- X > 0, !, fac(X - 1, Y1), Y is Y1 * X. fac(X, Y) :- X =:= 0, Y is 1.
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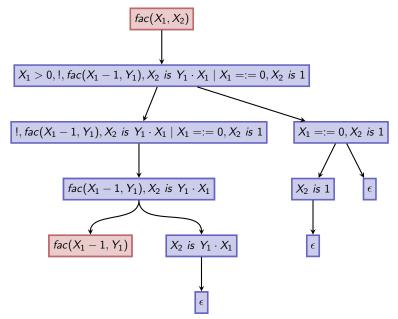
From Programs To Graphs - Nonterminating Construction



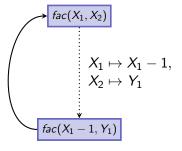
From Programs To Graphs - Split rule



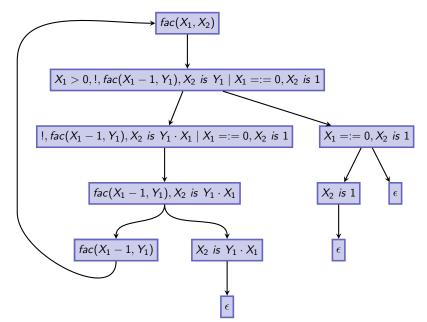
From Programs To Graphs - Split rule



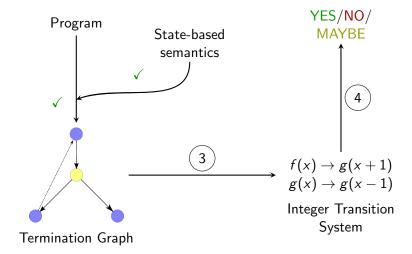
From Programs To Graphs - Instance Rule



From Programs To Graphs - Final Result



Roadmap



Integer Transition Systems

Integer Transition System:

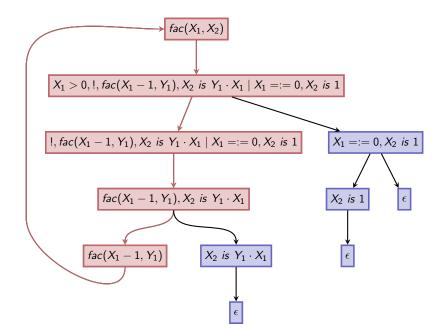
$$f(x) \rightarrow f(x+1)$$
 | $x < 0$

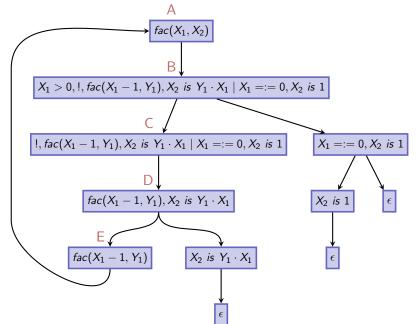
Given: Some Termination Graph

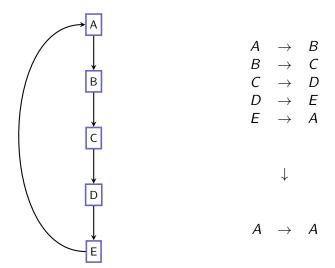
Goal: Integer Transition System that terminates if all runs described by the Termination Graph terminate

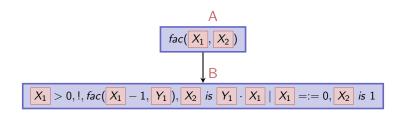
Idea: Encode graph locally, node by node

Paths in Graph \approx Evaluations

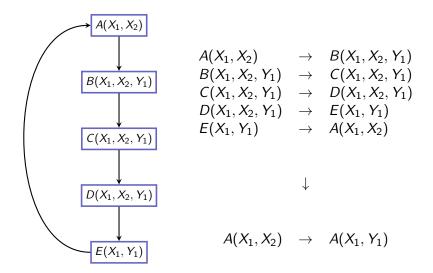




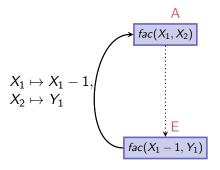


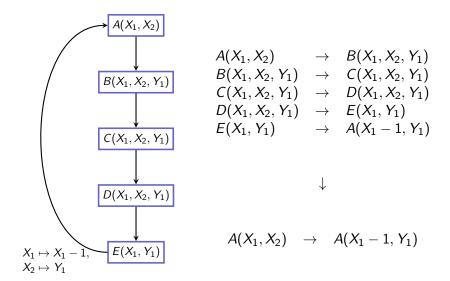


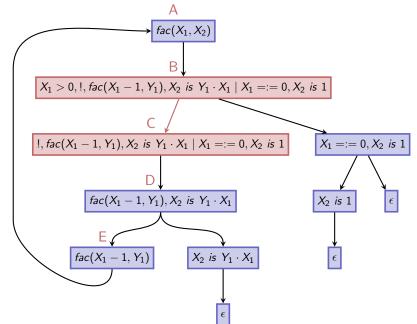
$$A(X_1, X_2) \rightarrow B(X_1, X_2, Y_1)$$

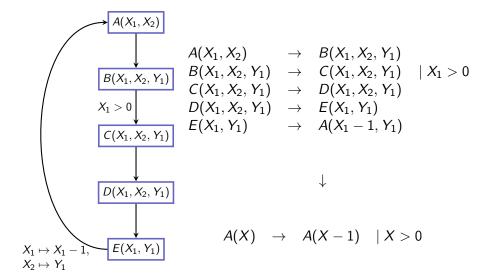


From Programs To Graphs - Instance Rule

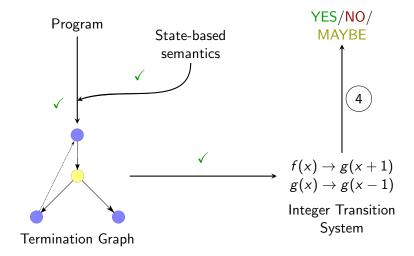








Roadmap

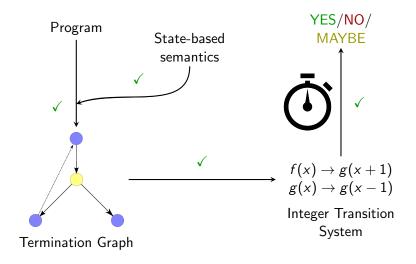


Termination of Integer Transition Systems

Well-studied problem

Use known techniques to show termination (Transition Invariants, Podelski and Rybalchenko, 2004)

Roadmap



Stopwatch: http://icons8.com/

Results - Compared Approaches

Term Rewriting with Argument Filter Directly from Program (Termination Proofs, Schneider-Kamp et al., 2009)

Dependency Triples Graph Construction (Dependency Triples, Ströder et al., 2011)

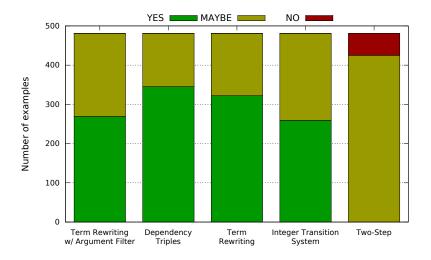
Term Rewriting Graph Construction (Symbolic Evaluation Graphs, Giesl et al., 2012)

Integer Transition Systems This approach

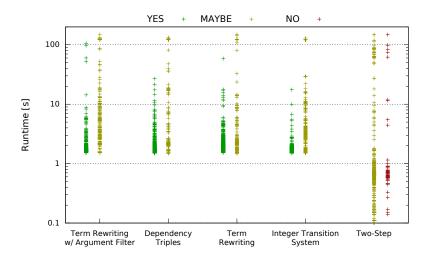
(Analysis of Arithmetic Prolog Programs, Weinert, 2015)

Two-Step Constraint Satisfaction Problem (Non-termination Analysis, Voets et al., 2011)

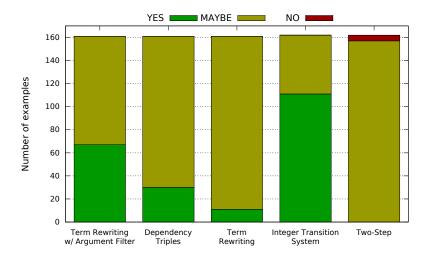
Results - Logic Benchmarks - Power



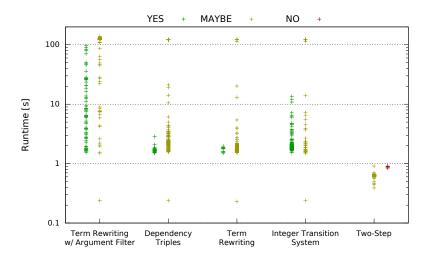
Results - Logic Benchmarks - Runtime



Results - Numerical Benchmarks - Power



Results - Numerical Benchmarks - Runtime



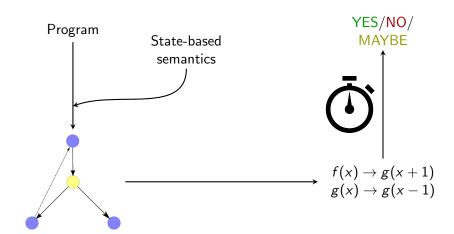
Contributions - Theoretical

- Extended construction of Termination Graphs, taking arithmetic comparisons and evaluations into account
- Developed new construction of Integer Transition System from Termination Graphs
- Extended abstract state to store arithmetic knowledge
- Separated abstract semantics and termination analysis
- Proved soundness of all steps of the construction

Contributions - Practical

- Implemented extension of construction of Termination Graphs, optimization through SMT solver
- Implemented construction of Integer Transition Systems from Termination Graphs
- Added 162 numerical benchmarks to benchmark suite
- Performed experiments comparing this approach to existing ones

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